

The Heptane Static Worst-Case Execution Time Estimation Tool

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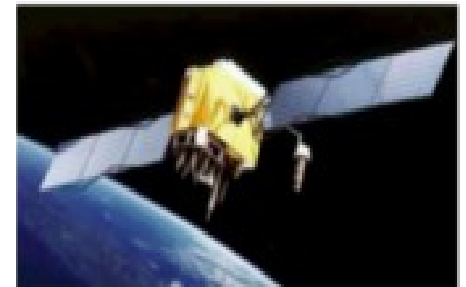
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Introduction



Real-time system
Timing guarantees
Worst-Case Execution Time (WCET)

- › Static methods
- › Measurement-based
- › Probabilistic



Heptane: a static analysis tool

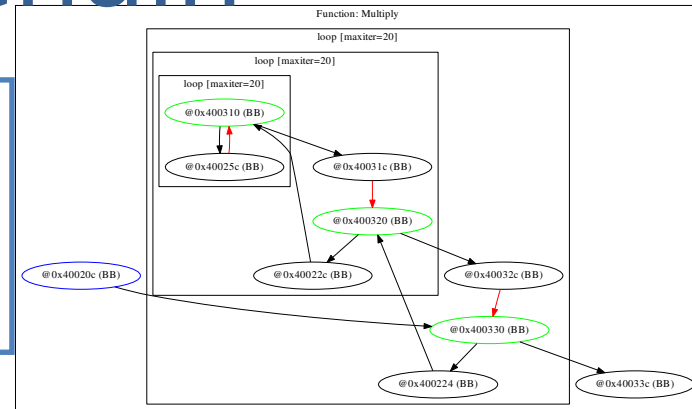
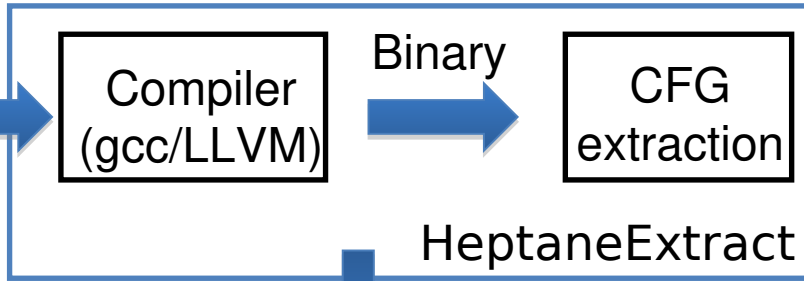
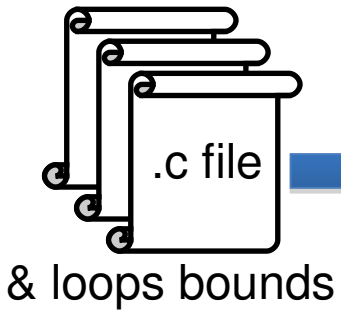
- Program structure analysis
- Flow information
- Binary level
- Model of the architecture

About HEPTANE

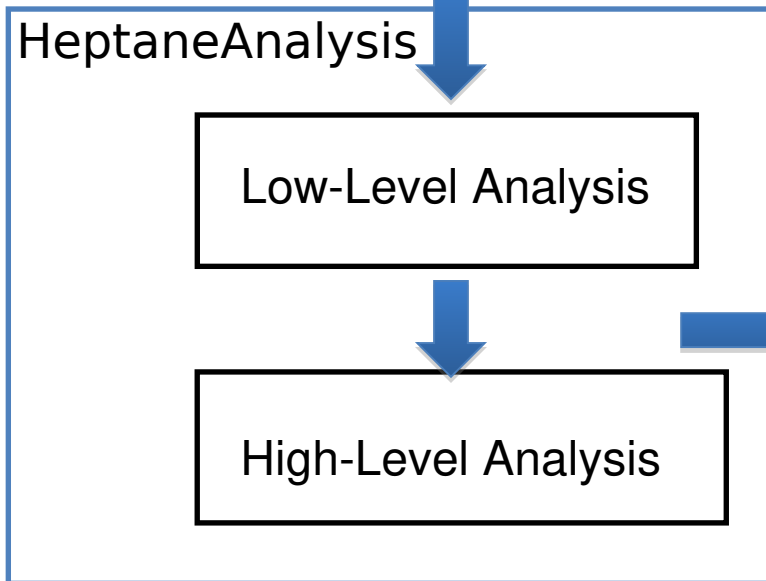
- Developed at IRISA/University of Rennes
 - Stable research prototype
 - Licenced under GPL
 - C++ code
 - Supported architectures
 - MIPS
 - Partially: ARM v7, Patmos (under construction)
- The first version has been designed around 2000
 - Many updates since then!
 - Available online

<https://team.inria.fr/pacap/software/heptane/>

Heptane toolchain



XML file



```

0           register int   Outer, Inner, Index;
0
1 20 21     for (Outer = 0; Outer < UPPERLIMIT; Outer++) {
0           ANNOT_MAXITER(20);
20 400 420  for (Inner = 0; Inner < UPPERLIMIT; Inner++) {
0           ANNOT_MAXITER(20);
  
```

- WCET
- Statistics
- HTML

How to use it

- Source code in C (alternatively binary)
 - With annotation for the loops bounds:
ANNOT_MAXITER(cst)

Include needed

```
#include <annot.h>
```

```
void main() {
```

```
    int i,a;
```

```
    a=0;
```

```
    for(i=0;i<10;i++){
```

```
        ANNOT_MAXITER(10);
```

```
        a++;
```

```
    }
```

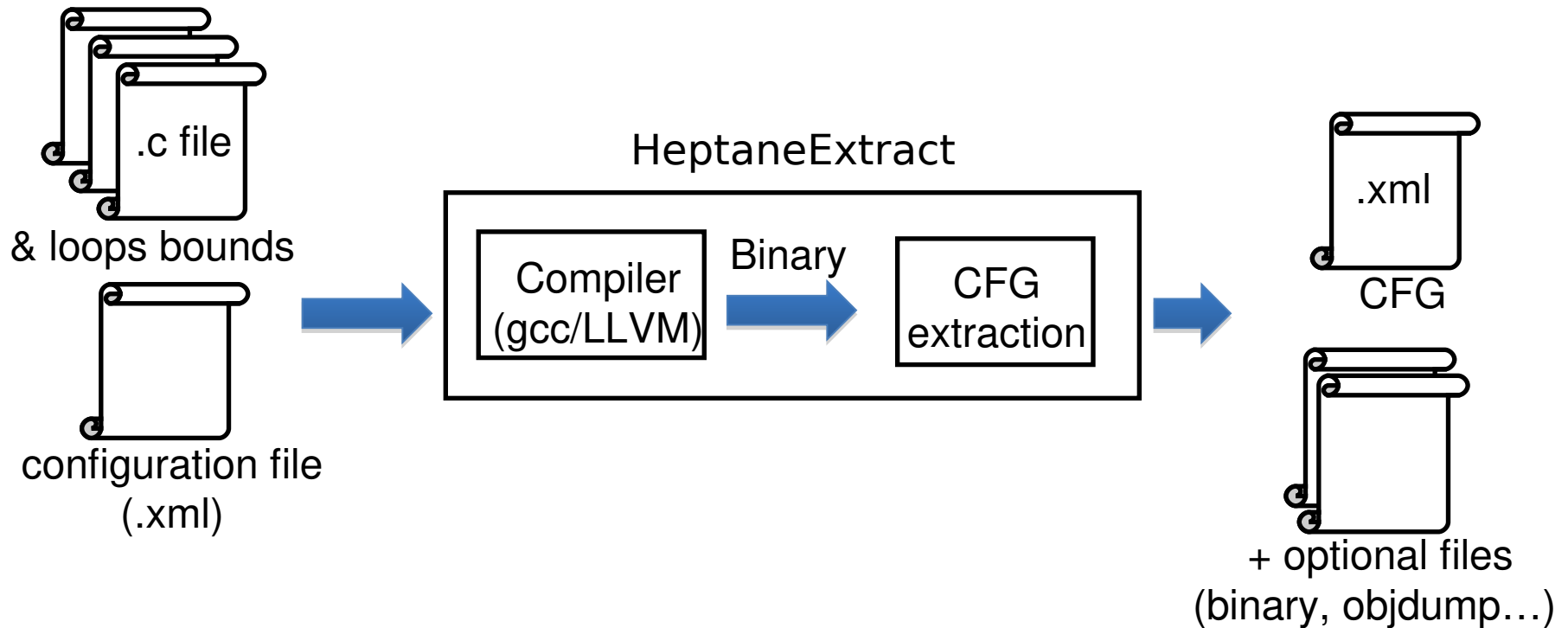
```
}
```

ANNOT_MAXITER for each loop

Some restrictions :

- No pointers
- No malloc
- No recursion
- Natural loops

How to use it: HeptaneExtract



CFG.xml overview

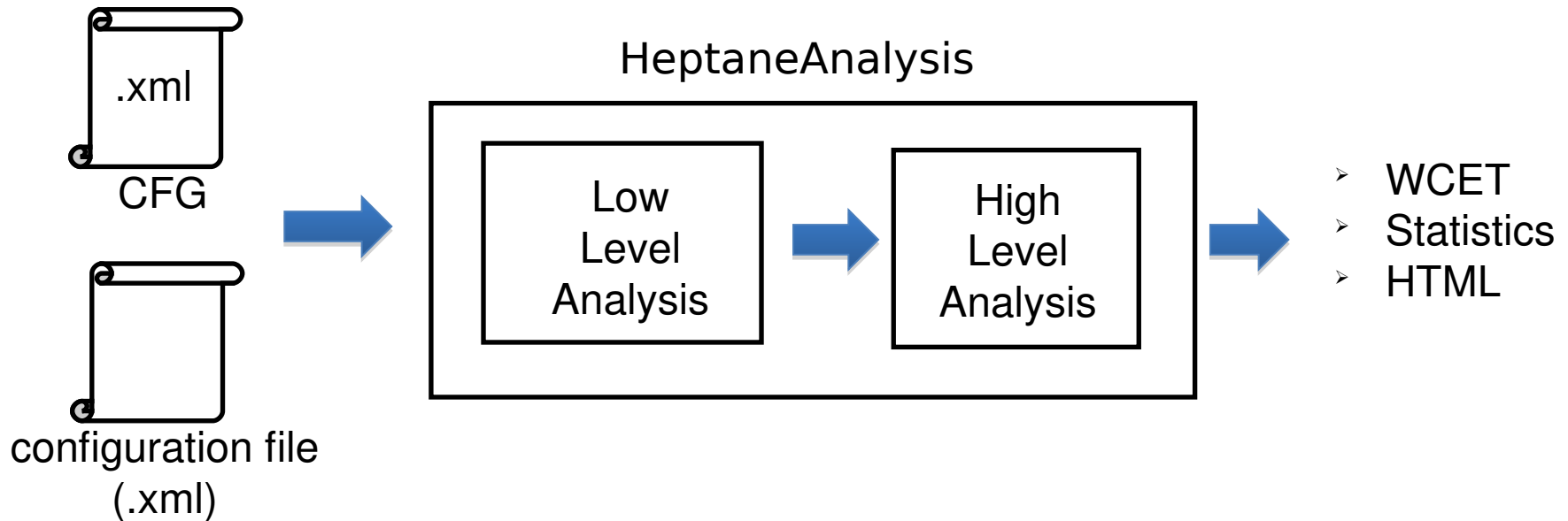
```

<!DOCTYPE PROGRAM SYSTEM 'cfglib.dtd'>
<PROGRAM id="1" name="simple" entry="0" >
<CFG id="0" name="main" startnode="2" endnodes="3, " >
<NODE id="2" type="BasicBlock" >
<INSTRUCTION id="4" asm_type="Code" code="addiu sp,sp,-8" >
<ATTRS_LIST>
<ATTR type="address" name="address">
<ACCESS type="read" seg="code" varname="" precision="1">
<ADDRSIZE begin="0x400018" size="4"/>
</ACCESS>
</ATTR>
</ATTRS_LIST>
</INSTRUCTION>
<INSTRUCTION id="5" asm_type="Code" code="sw zero,4(sp)" >
<ATTRS_LIST>
<ATTR type="address" name="address">
<ACCESS type="read" seg="code" varname="" precision="1">
<ADDRSIZE begin="0x40001c" size="4"/>
</ACCESS>
</ATTR>
</ATTRS_LIST>
</INSTRUCTION>
<EDGE id="25" origin="2" destination="17" >
</EDGE>
<EDGE id="26" origin="8" destination="17" >
</EDGE>
<EDGE id="27" origin="17" destination="8" >
</EDGE>
<EDGE id="28" origin="17" destination="3" >
</EDGE>
<LOOP id="29" head="17" nodes="17, 8, " backedges="26, " >
<ATTRS_LIST>
<ATTR type="integer" name="maxiter" value="10" />
</ATTRS_LIST>
</LOOP>

```

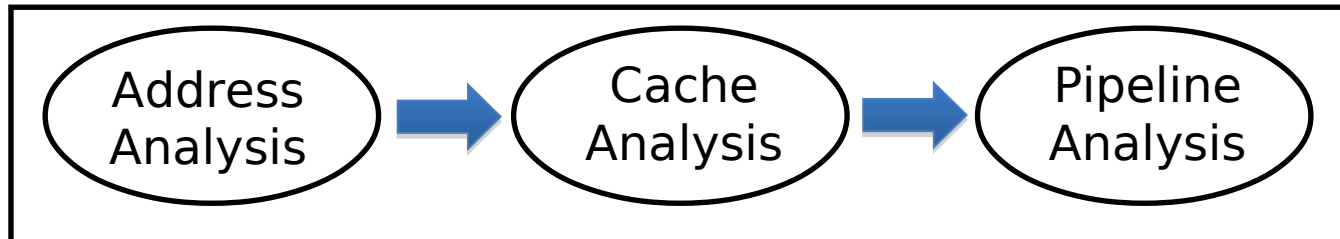
- 1 cfg per function
- Instruction listing filling basic-blocks
- Instruction addresses
- Edges between basic-blocks
- Attributes (e.g. loop bounds)

How to use it: HeptaneAnalysis



Low-level analysis

- Account for micro-architecture
- Timing information per basic-block
- Enrich the CFG.xml



Data address analysis

- Conservatively determines addresses of referenced data
- Instruction addresses from the binary
- Stack analysis per function
 - first instruction of functions
 - calling context
- Basic block granularity
 - for each access (load/store)
 - evaluate registers content + memory (new)

Cache analysis

- › Instruction & data set-associative caches
 - › contextual analysis
 - › abstract interpretation (must, may, persistence)
- › Multi-level non inclusive cache hierarchies
 - › write-through, no-write allocate data caches
 - › LRU, PLRU, FIFO, MRU, Random
- › Cache Hit Miss Classification (CHMC)
 - › Always hit, Always miss, First miss, Not-classified

Pipeline analysis

- Compute the timing of basic blocks
 - simple in-order pipeline
 - free of timing anomalies
 - fetch/memory stage depends on cache analysis

High-level analysis

- Determine critical path
- Context sensitive (calling context)
- Implicit Path Enumeration Technique (IPET)

maximize $\sum_{i \in BasicBlocks} T_i \times f_i$

T_i → execution time
constant computed by low-level analysis

f_i → # executions

Install, use and hack

Install

GNU/Linux & MacOS

download sources from : <https://team.inria.fr/pacap/software/heptane/>

modify *install.sh* to your configuration

dependencies: libxml2, ILP solver (lp-solve, cplex)

Use

create a config file (sample in *config_files/* folder)

launch *HeptaneExtract*

launch *HeptaneAnalysis*

collect results

Hack

architecture specificities in *src/Common/ArchitectureDependent/*
analysis in *src/HeptaneAnalysis/src/Specific/*

Research based on Heptane

- Cache analysis (cache hierarchies, multi-cores)

D. Hardy, I. Puaut. WCET analysis of instruction cache hierarchies. Journal of Systems Architecture, volume 57, issue 7, pages 677-694, August 2011.

- Branch prediction (analysis, compiler-directed)

A. Colin, I. Puaut. Worst Case Execution Time Analysis for a Processor with Branch Prediction. Real-Time Systems, April 2000.

- Cache locking, partitioning

I. Puaut. Cache Analysis vs Static Cache Locking for Schedulability Analysis in Multitasking Real-Time Systems. WCET Workshop, ECRTS 2002.

- Scratchpad management

J-F. Deverge, I. Puaut. WCET-directed dynamic scratchpad memory allocation of data. ECRTS 2007.

- CRPD estimation, Static probabilistic WCET, Traceability of flow information, ...

<https://team.inria.fr/pacap/members/isabelle-puaut/>

Conclusion

Heptane WCET analysis tool

- open source

- estimate WCET from binary

- advanced in static cache analysis

- Side product : CFG management library

Future work

- automatic loop bound extraction

- further improve address analysis

- increase number of supported architecture

Any help is welcome

Questions?



<https://team.inria.fr/pacap/software/heptane/>

Specific questions, skype: isabellepuaut